

ML2R Coding Nuggets

Solving Linear Programming Problems

Pascal Welke*
Machine Learning Rhine-Ruhr
University of Bonn
Bonn, Germany

Christian Bauckhage†
Machine Learning Rhine-Ruhr
Fraunhofer IAIS
St. Augustin, Germany

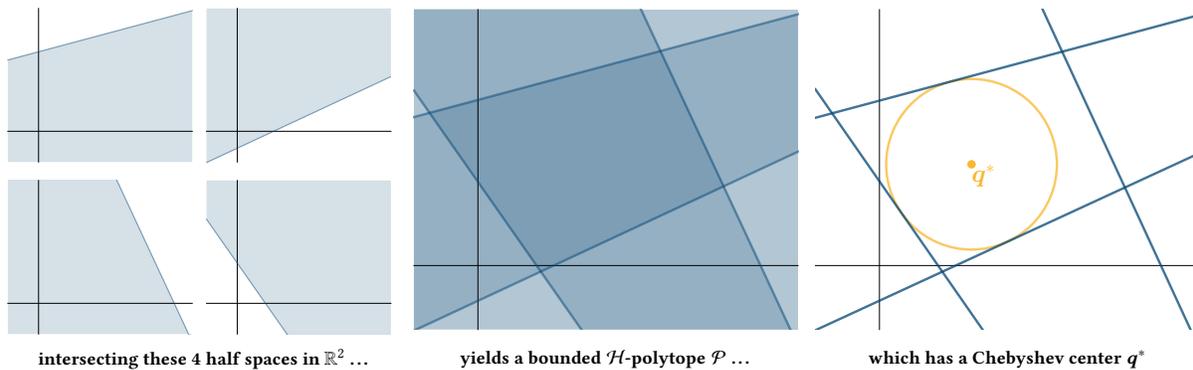


Figure 1: The Chebyshev center q^* of a bounded convex \mathcal{H} -polytope \mathcal{P} is the center point of its largest inscribed ball.

ABSTRACT

This note discusses how to solve linear programming problems with *SciPy*. As a practical use case, we consider the task of computing the Chebyshev center of a bounded convex polytope.

1 INTRODUCTION

Many optimization problems in data science, machine learning, and artificial intelligence are **linear programming** problems. In general, these are written as

$$\begin{aligned} z^* &= \operatorname{argmin}_{z \in \mathbb{R}^n} c^\top z \\ \text{s.t. } & Az \leq b \\ & Cz = d \end{aligned} \quad (1)$$

and the crux of the matter is perfectly summarized by Boyd and Vandenberghe [5, chapter 1]: “There is no simple analytical formula for the solution of a linear program ... but there are a variety of very effective methods for solving them, including Dantzig’s simplex method, and the more recent interior-point methods ... We can easily solve problems with hundreds of variables and thousands of constraints on a small desktop computer, in a matter of seconds.”

Indeed, *SciPy* has us covered. Its `optimize` package provides the function `linprog` which implements simplex- and interior-point solvers for problems of the form in (1). While the use of `linprog` is straightforward, it can be challenging to rewrite a given problem

such that it fits the function’s input requirements. This is what this note is all about.

As a practical example, we consider the problem of estimating the Chebyshev center of a bounded nonempty convex polytope (see Fig. 1). First, we briefly look at the underlying theory (section 2) and then solve the corresponding linear program using `linprog` (section 3).

Readers who would like to experiment with our code should be passingly familiar with *NumPy* and *SciPy* [11] and only need to

```
import numpy as np
import scipy.optimize as opt
```

2 THEORY

Next, we first clarify basic terms and definitions and then show that the problem of computing Chebyshev centers can be seen as a special case of (1).

Terms and Definitions

Throughout, we understand the **Chebyshev center** of a bounded nonempty convex \mathcal{H} -polytope to be the center point of the largest Euclidean ball fully inscribed within said polytope.¹ An illustration of this meaning is shown in Fig. 1.

* 0000-0002-2123-3781
† 0000-0001-6615-2128

¹Unfortunately, the term is also used to describe closely related yet different geometric entities: en.wikipedia.org/wiki/Chebyshev_center.

Moreover, following Ziegler [13], we say that an m -dimensional \mathcal{H} -polytope is a set

$$\mathcal{P} = \left\{ \mathbf{x} \in \mathbb{R}^m \mid \mathbf{W}\mathbf{x} \leq \boldsymbol{\theta} \right\} \quad (2)$$

where the matrix-vector expression $\mathbf{W}\mathbf{x} \leq \boldsymbol{\theta}$ is a convenient shorthand for a collection of $i = 1, \dots, p$ inequalities $\mathbf{w}_i^\top \mathbf{x} \leq \theta_i$.

In other words, each point \mathbf{x} inside of an \mathcal{H} -polytope \mathcal{P} is a solution to a system of p linear inequalities. Since each of these inequalities defines a half-space in \mathbb{R}^m , we recognize the definition in (2) to describe a convex set, namely an intersection of finitely many half-spaces, hence the name \mathcal{H} -polytope. Finally, if there exists some number $R \in \mathbb{R}$ such that $\|\mathbf{x}\| \leq R$ for all $\mathbf{x} \in \mathcal{P}$, then \mathcal{P} is bounded.

An m -dimensional Euclidean ball is yet another, arguably more prosaic bounded convex set, namely

$$\mathcal{B} = \left\{ \mathbf{x} \in \mathbb{R}^m \mid \|\mathbf{x} - \mathbf{q}\| \leq r \right\}. \quad (3)$$

Here, the two parameters $\mathbf{q} \in \mathbb{R}^m$ and $r \in \mathbb{R}$ simply characterize the ball's center point and radius.

Chebyshev Centers and Linear Programming

Since a Euclidean ball is defined in terms of its center and radius, to estimate the Chebyshev center of a bounded \mathcal{H} -polytope \mathcal{P} is to estimate the parameters \mathbf{q}^* and r^* of the largest ball inscribed in it.

Without retracing all their arguments, we simply note that Boyd and Vandenberghe [5] show that these parameters coincide with the solution to the following constrained optimization problem

$$\begin{aligned} \mathbf{q}^*, r^* = \operatorname{argmax}_{\mathbf{q}, r} \quad & r \\ \text{s.t.} \quad & r \cdot \|\mathbf{w}_i\| + \mathbf{w}_i^\top \mathbf{q} \leq \theta_i, \quad i = 1, \dots, p. \end{aligned} \quad (4)$$

Now, in order to see that (4) is indeed but a special case of (1), we proceed as follows:

- (1) we combine the two problem variables $\mathbf{q} \in \mathbb{R}^m$ and $r \in \mathbb{R}$ into a single vector $\mathbf{z} \in \mathbb{R}^{1+m}$ such that

$$\mathbf{z} = \begin{bmatrix} r \\ \mathbf{q} \end{bmatrix} \quad (5)$$

- (2) we introduce yet another vector $\mathbf{c} \in \mathbb{R}^{1+m}$, namely

$$\mathbf{c} = \begin{bmatrix} -1 \\ 0 \\ \vdots \\ 0 \end{bmatrix} \quad (6)$$

- (3) these vectors allow us to rewrite the maximization objective in equation (4) in terms of a minimization objective which involves an inner product, because

$$\operatorname{argmax}_{\mathbf{q}, r} r \Leftrightarrow \operatorname{argmin}_{\mathbf{q}, r} -r \Leftrightarrow \operatorname{argmin}_{\mathbf{z}} \mathbf{c}^\top \mathbf{z} \quad (7)$$

- (4) we introduce a matrix $\mathbf{A} \in \mathbb{R}^{p \times (1+m)}$ together with a vector $\mathbf{b} \in \mathbb{R}^p$ which are given by

$$\mathbf{A} = \begin{bmatrix} \|\mathbf{w}_1\| & \mathbf{w}_1^\top \\ \vdots & \vdots \\ \|\mathbf{w}_p\| & \mathbf{w}_p^\top \end{bmatrix} \quad \text{and} \quad \mathbf{b} = \begin{bmatrix} \theta_1 \\ \vdots \\ \theta_p \end{bmatrix} \quad (8)$$

- (5) together with \mathbf{z} , \mathbf{A} and \mathbf{b} allow us to write the collection of p individual inequality constraints in (4) in terms of a single matrix-vector expression, namely

$$\mathbf{A}\mathbf{z} \leq \mathbf{b} \quad (9)$$

All in all, our (re)definitions therefore allow us to (re)write the optimization problem in (4) as

$$\begin{aligned} \mathbf{z}^* = \operatorname{argmin}_{\mathbf{z} \in \mathbb{R}^{m+1}} \quad & \mathbf{c}^\top \mathbf{z} \\ \text{s.t.} \quad & \mathbf{A}\mathbf{z} \leq \mathbf{b}. \end{aligned} \quad (10)$$

This is now easily recognizable as a special case of the more general linear program in (1), namely as a linear program without equality constraints.

3 PRACTICAL COMPUTATION

In this section, we discuss how to solve the Chebyshev center problem using *SciPy*. Given what we just worked out (equations (5)–(10)), there is actually not much left to discuss except for details of the code in Listing 1.

In order to work with a specific (numerical) example, we resort to the $m = 2$ dimensional polytope \mathcal{P} in Fig. 1. This particular polytope results from plunging

$$\mathbf{W} = \begin{bmatrix} -0.26 & 0.97 \\ 0.42 & -0.91 \\ 0.91 & 0.42 \\ -0.82 & -0.57 \end{bmatrix} \quad \text{and} \quad \boldsymbol{\theta} = \begin{bmatrix} 5.0 \\ 1.0 \\ 8.0 \\ -1.5 \end{bmatrix} \quad (11)$$

into (2). Implementing both, matrix \mathbf{W} and vector $\boldsymbol{\theta}$, in terms of *NumPy* arrays is straightforward

```
matW = np.array([[ -0.26,  0.42,  0.91, -0.82],
                 [ 0.97, -0.91,  0.42, -0.57]]).T
vecT = np.array([5.0, 1.0, 8.0, -1.5])
```

Likewise, it is easy to compute the norms $\|\mathbf{w}_i\|$ of the rows of matrix \mathbf{W} . We simply use the following recipe (for an explanation of the rationale behind this one-liner, see [1])

```
rowNrmW = np.sqrt(np.sum(matW**2, axis=1))
```

Next, we need to set up matrix \mathbf{A} and vector \mathbf{b} as introduced in equation (8). To accomplish this, we may proceed as follows

```
matA = np.vstack((rowNrmW, matW.T)).T
vecB = vecT
```

The last ingredient of our linear program is the vector \mathbf{c} in (6). Since our current problem is $1 + m = 3$ dimensional, we can simply instantiate it like so

```
vecC = np.array([-1, 0, 0])
```

or, more generally, write

```
vecC = np.zeros(matA.shape[1])
vecC[0] = -1
```

At this point, we are good to go and can invoke function `linprog` in *SciPy*'s `optimize` module. It is called with several parameters of which the following are most relevant in our current setting:

- \mathbf{c} is a 1D array representing the coefficient vector \mathbf{c} of the linear objective function we wish to solve
- $\mathbf{A_ub}$ and $\mathbf{b_ub}$ are a 2D and a 1D array which represent the matrix \mathbf{A} and vector \mathbf{b} which define the inequality (or *upper bound*) constraints of our problem

Solving Linear Programming Problems

- although we do not actually need them here, we should mention that the parameters `A_eq` and `b_eq` are a 2D and a 1D array which would represent matrix C and vector d which occur in (1) and specify potential equality constraints
- finally, the parameter `method` is a string used to indicate which solver we wish to apply; its default value amounts to 'interior-point', another reasonable choice would be 'revised simplex'.

Hence, for our current problem, we may use `linprog` as follows

```
result = opt.linprog(vecC, A_ub=matA, b_ub=vecB)
```

This will cause `result` to be a variable of type `OptimizeResult` which is a class that `SciPy` uses in order to summarize the outcome of an optimization procedure. To have a look at our result, we may therefore simply

```
print (result)
```

which provides us with the following status report

```
con: array([], dtype=float64)
fun: -2.666293247661684
message: 'Optimization terminated successfully.'
nit: 3
slack: array([0., 0., 1.38093669, 0.])
status: 0
success: True
x: array([2.66629325, 2.87626447, 3.16518324])
```

Its content is largely self explanatory; explanations of the more cryptic entries can be found in the `SciPy` documentation. The most important thing to us is that the field `x` contains the solution z^* to our problem. In other words, something like

```
vecZ = result.x
r, c = vecZ[0], vecZ[1:]
```

will provide us with radius `r` and center point `c` of the largest ball inscribed within our polytope.

3.1 A Note on “Corner Cases”

Intuitively, it is clear that we can only find a Chebyshev-center of an *nonempty* polytope P . However, no-one keeps you from defining an empty polytope by adding the inequality $[1 \ 0] x \leq -3$:

```
matW = np.array([[[-0.26, 0.42, 0.91, -0.82, 1],
                  [0.97, -0.91, 0.42, -0.57, 0]]]).T
vecT = np.array([5.0, 1.0, 8.0, -1.5, -3])
```

There is no $x \in \mathbb{R}^2$ that satisfies all constraints at once. Luckily, `linprog` notices this and results in the following output:

```
con: array([], dtype=float64)
fun: -0.6671837546213493
message: 'The algorithm terminated successfully and determined
         that the problem is infeasible.'
nit: 4
slack: array([3.47102113, 1.09568071, 6.71856309,
              -1.43242614, -3.90269694])
status: 2
success: False
x: array([0.66718375, 0.23551318, 0.94865877])
```

Be careful! As you can see, the result tells you that no solution exists, but at the same time happily gives you a “solution” vector `result.x`. Hence, you *must always* check the value of `result.success` before using `result.x`!

Listing 1: solving the linear program in (10)

```
matW = np.array([[[-0.26, 0.42, 0.91, -0.82],
                  [0.97, -0.91, 0.42, -0.57]]]).T
vecT = np.array([5.0, 1.0, 8.0, -1.5])

rowNrmW = np.sqrt(np.sum(matW**2, axis=1))

matA = np.vstack((rowNrmW, matW.T)).T
vecB = vecT

vecC = np.zeros(matA.shape[1])
vecC[0] = -1

result = opt.linprog(vecC, A_ub=matA, b_ub=vecB)
```

4 SUMMARY AND OUTLOOK

This short note discussed how to solve linear programming problems using the method `linprog` contained in `SciPy`'s `optimize` package. We saw that the use of `linprog` is intuitive, as long as we can express the problem we are dealing with in the form expected by the function. To see how to bring a given problem into this particular form, we considered the Chebyshev center problem and rewrote the underlying linear program correspondingly.

While our practical example of computing the Chebyshev center of a polytope, or, equivalently, of computing its largest inscribed ball was chosen for didactic purposes, we note that Euclidean balls are a staple of intelligent data analysis [2–4, 8, 10, 12]. Moreover and more recently, Euclidean balls have also been used successfully for informed or structured representation learning [6, 7, 9] and we will return to these ideas in later notes.

ACKNOWLEDGMENTS

This material was produced within the Competence Center for Machine Learning Rhine-Ruhr (ML2R) which is funded by the Federal Ministry of Education and Research of Germany (grant no. 01IS18038A). The authors gratefully acknowledge this support.

REFERENCES

- [1] C. Bauckhage. 2015. NumPy / SciPy Recipes for Data Science: Computing Nearest Neighbors. researchgate.net. [dx.doi.org/10.13140/RG.2.1.4602.0564](https://doi.org/10.13140/RG.2.1.4602.0564).
- [2] C. Bauckhage, M. Bortz, and R. Sifa. 2020. Shells within Minimum Enclosing Balls. In *Proc. DSAA*. IEEE.
- [3] C. Bauckhage, R. Sifa, and T. Dong. 2019. Prototypes within Minimum Enclosing Balls. In *Proc. ICANN*.
- [4] A. Ben-Hur, D. Horn, H.T. Siegelmann, and V. Vapnik. 2001. Support Vector Clustering. *J. of Machine Learning Research* 2 (2001).
- [5] S. Boyd and L. Vandenberghe. 2004. *Convex Optimization*. Cambridge University Press.
- [6] T. Dong, C. Bauckhage, H. Jin, J. Li, O. Cremers, D. Speicher, A.B. Cremers, and J. Zimmermann. 2019. Imposing Category Trees onto Word-Embeddings Using a Geometric Construction. In *Proc. ICLR*.
- [7] T. Dong, Z. Wang, J. Li, C. Bauckhage, and A.B. Cremers. 2019. Triple Classification Using Regions and Fine-Grained Entity Typing. In *Proc. AAAI*.
- [8] G.D. Evangelidis and C. Bauckhage. 2013. Efficient Subframe Video Alignment Using Short Descriptors. *IEEE Trans. Pattern Analysis and Machine Intelligence* 35, 10 (2013).
- [9] T. Le, H. Vu, T.D. Nguyen, and D. Phung. 2018. Geometric Enclosing Networks. In *Proc. IJCAI*.
- [10] J. Lee and D. Lee. 2005. An Improved Cluster Labeling Method for Support Vector Clustering. *IEEE Trans. Pattern Analysis and Machine Intelligence* 27, 3 (2005).
- [11] T.E. Oliphant. 2007. Python for Scientific Computing. *Computing in Science & Engineering* 9, 3 (2007).
- [12] D.M.J. Tax and R.P.W. Duin. 2004. Support Vector Data Description. *Machine Learning* 54, 1 (2004).
- [13] G.M. Ziegler. 1995. *Lectures on Polytopes*. Springer.